

```
procedure INIT-PDA
{ Invoked when the router comes up. }
begin
    Initialize all tables;
    call PDA;
end INIT-PDA

procedure PDA
{ Executed at each router i. Invoked when an event occurs }
begin
    (1) call NTU;
    (2) call MTU; /* Updates  $T^i$  */
    (3) if (there are changes to  $T^i$ ) then
        Compose an LSU message consisting of topology
        differences using add, delete
        and change link entries;
    end if
    (4) Within a finite amount time, send the
        LSU message to all neighbors;
end PDA
```

FIG. 1

```

procedure NTU
begin
  (1) if (LSU message is received from a neighbor  $k$ ) then
    (1a) Update neighbor table  $T_k^i$ . That is, add links,
        delete links or change links according to the
        specification of each entry in the LSU;
    (1b) Run Dijkstra's shortest path algorithm
        on the resulting topology  $T_k^i$ ; /*This results in
        finding minimum distances from  $k$  to all other
        nodes in  $T_k^i$ . Note  $T_k^i$  is a tree*/
    (1c) Update  $D_{jk}^i$  with new distances in  $T_k^i$ ;
  end if
  (2) if (adjacent link  $(i, k)$  is up) then
    Update  $l_k^i$  and send an LSU message to the
    neighbor  $k$  with link information of all links in
    its main topology table  $T^i$ ;
  endif
  (3) if (cost of an adjacent link  $(i, k)$  changed)then
    Update  $l_k^i$ ;
  endif
  (4) if (adjacent link  $(i, k)$  failed)then
    Update  $l_k^i$  and clear the table  $T_k^i$ ;
  endif
end NTU

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FIG. 2

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procedure MTU at router  $i$ 
begin
  (1)  $oldT^i \leftarrow T^i$ ; /* Save copy */
  (2) if (node  $j$  occurs in at least one of  $T_k^i$ ) then
    add  $j$  to the main topology table  $T^i$ 
  end if
  (3) for each node  $j$  in  $T^i$  do
     $MIN \leftarrow \min \{D_{jk}^i + l_k^i \mid k \in N^i\}$ ;
    let  $p$  be such that  $MIN = (D_{jp}^i + l_p^i)$ ;
    /* Neighbor  $p$  is the preferred neighbor for
    destination  $j$ . Ties are broken in favor of
    lower address neighbor */
  done
  (4) for each  $j$  in  $T^i$  and its preferred neighbor  $p$  do
    Copy all links  $(j, n)$  from  $T_p^i$  to  $T^i$ ;
    /* i.e., copy all links in  $T_p^i$  for which
     $j$  is the head node */
  done
  (5) Update  $T^i$  with information of each  $l_k^i$ ;
  (6) Run Dijkstra's shortest path algorithm on  $T^i$ 
    and remove those links in  $T^i$  that are not
    part of the shortest path tree;
  (7) Update  $D_j^i$  with new distances in  $T^i$ ;
  (8) Compare  $oldT^i$  with  $T^i$  and note all differences;
end MTU

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FIG. 3

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procedure MPDA at router  $i$ 
{invoked when an event occurs}
begin
  (1) call NTU;
  (2) if (node is in PASSIVE state) then
    (2a) call MTU; /* update  $T_j^i$  and  $D_j^i$  */
    (2b)  $F D_j^i \leftarrow \min\{F D_j^i, D_j^i\}$ ;
  endif
  (3) if (node is in ACTIVE state and the
        last ACK is received) then
    (3a)  $temp_j^i \leftarrow D_j^i$ ; set node to PASSIVE state;
    (3b) call MTU to update  $T_j^i$ ;
    (3c)  $F D_j^i \leftarrow \min\{temp_j^i, D_j^i\}$ 
  endif
  (4)  $S_j^i \leftarrow \{k | D_j^i < F D_j^i\}$ ;
  (5) if (changes occur in  $T_j^i$ ) then
    Set node to ACTIVE state;
  endif
  if (no changes occur in  $T_j^i$  and the event is
      the last ACK) then
    Set node to PASSIVE state;
  endif
  (6) if (there are changes to  $T_j^i$ ) then
    Compose anew LSU with the topology
    changes expressed as add link,
    delete link and change link;
  end if
  (7) if (input event received is an LSU message) then
    Add the ACK entry to newly composed LSU;
  endif
  (8) Send the new LSU message.
end MPDA

```

FIG. 4

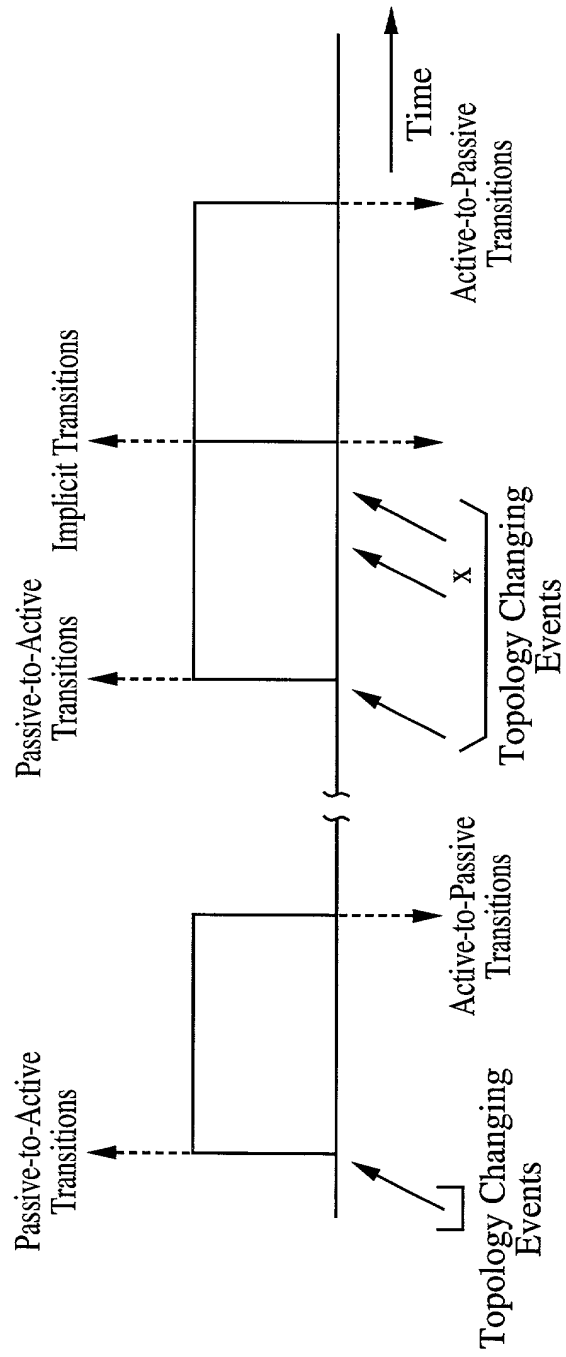


FIG. 5

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Procedure IH
begin
  (1)  $\forall k \notin S_j^i: \phi_{jk}^i \leftarrow 0;$ 
  (2) if  $(|S_j^i| = 1)$  then
     $\forall k \in S_j^i: \phi_{jk}^i \leftarrow 1;$ 
  endif
  (3) if  $(|S_j^i| > 1)$  then
    
$$1 - \frac{D_{jk}^i + l_k^i}{\sum_{m \in S_j^i} (D_j^i + l_m^i)}, \quad \forall k \in S_j^i;$$

  endif
end IH

```

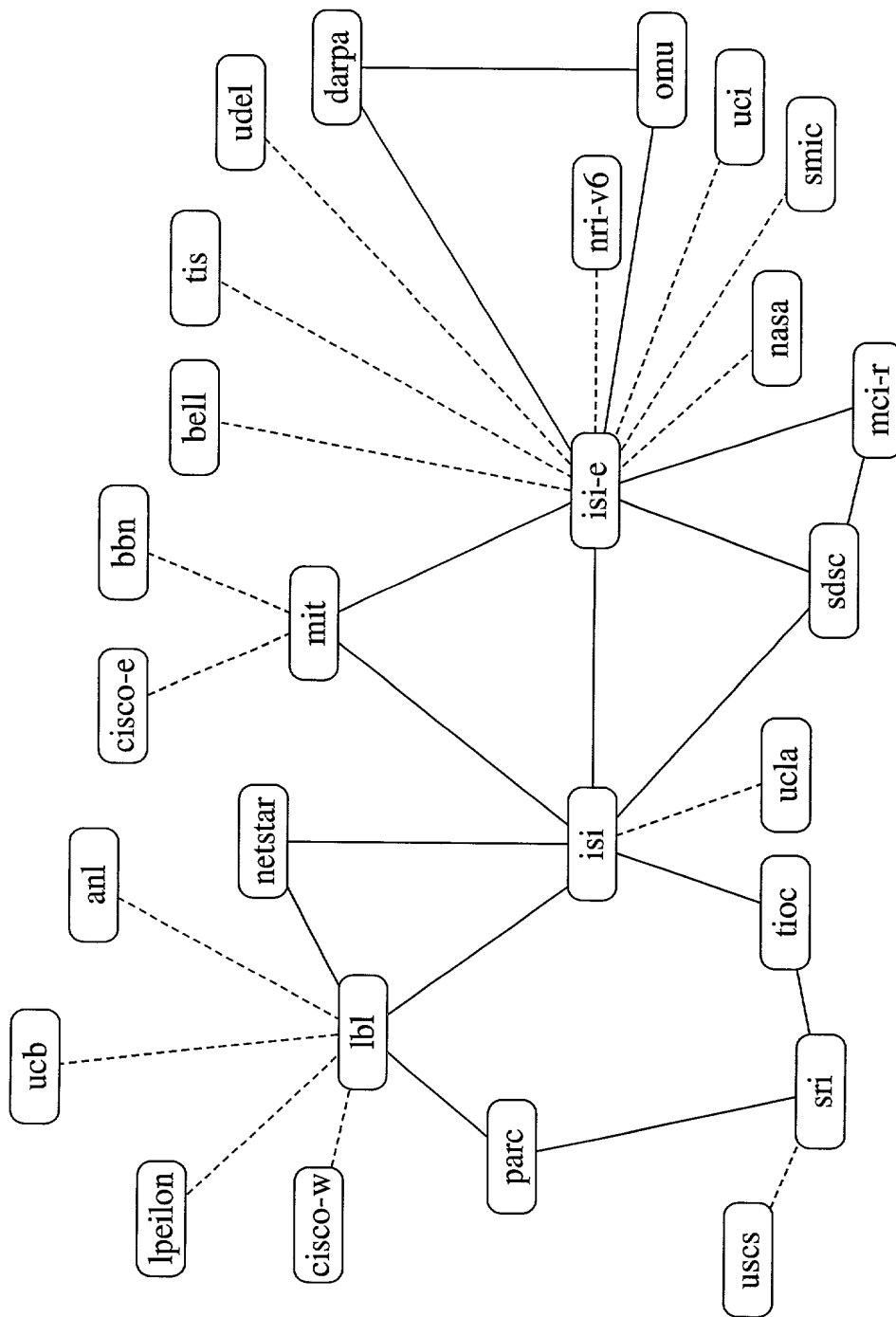
FIG. 6

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Procedure AH
begin
  (1)  $D_{min}^{ij} \leftarrow \min \{D_{jk}^i + l_k^i | k \in S_j\}$ ;
  (2) let  $D_{min}^{ij} = (D_{jk_0}^i + l_{k_0}^i)$ 
      // that is,  $k_0$  be the neighbor
      that offers the minimum
  (3) foreach  $k \in S_j^i$  do
       $a_{jk}^i \leftarrow D_{jk}^i + l_k^i - D_{min}^{ij}$ ;
    done
  (4)  $\Delta \leftarrow \frac{1}{2} \min \{ \frac{\phi_{jk}^i}{a_{jk}^i} | k \in S_j^i \wedge a_{jk}^i \neq 0 \}$ ;
  (5) foreach  $k \neq k_0 \wedge k \in S_j^i$  do
       $\phi_{jk}^i \leftarrow \phi_{jk}^i - \Delta \times a_{jk}^i$ ;
    done
  (6) foreach  $k = k_0$  do
       $\phi_{jk}^i \leftarrow \phi_{jk}^i + \sum_{q \in S_{jk}^i} \Delta \times a_{jq}^i$ ;
    done
end AH

```

FIG. 7



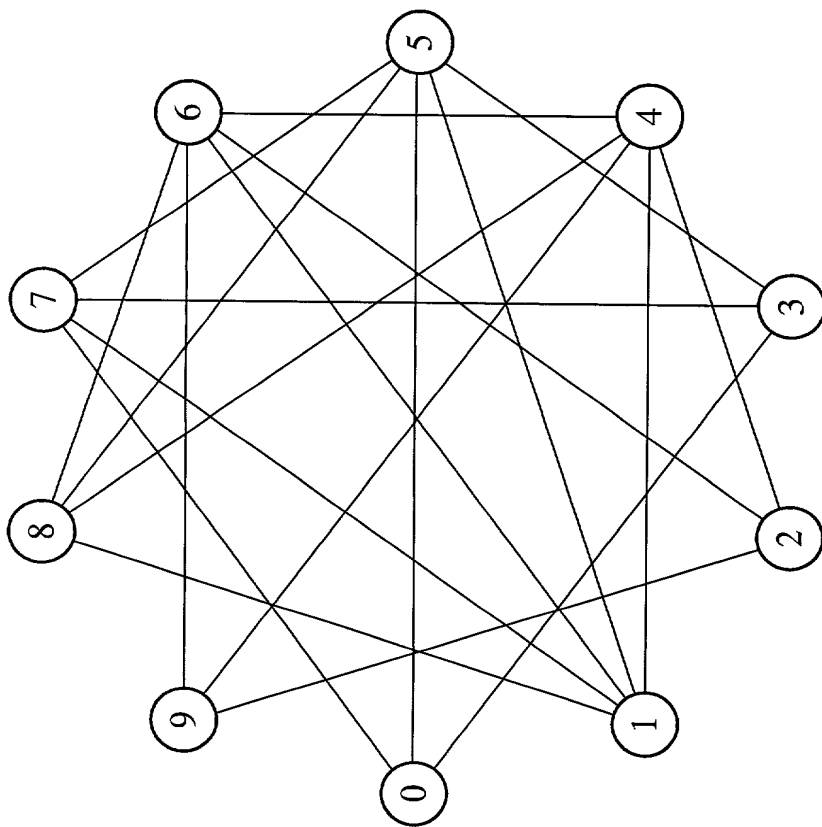


FIG. 9

FIG. 10 is a line graph showing Average Delays in Milliseconds versus Flow IDs for three different optimization methods: 'OPT', 'MP-TL-10-TS-2', and 'OPT-25'. The graph illustrates the performance of these methods across 10 flow IDs. The Y-axis represents Average Delays in Milliseconds, ranging from 1 to 4. The X-axis represents Flow IDs, ranging from 1 to 10. The 'OPT' method (solid line with diamond markers) generally shows the lowest delays, while the 'MP-TL-10-TS-2' method (dashed line with square markers) and 'OPT-25' method (dotted line with square markers) show higher delays, with 'OPT-25' often being the highest.

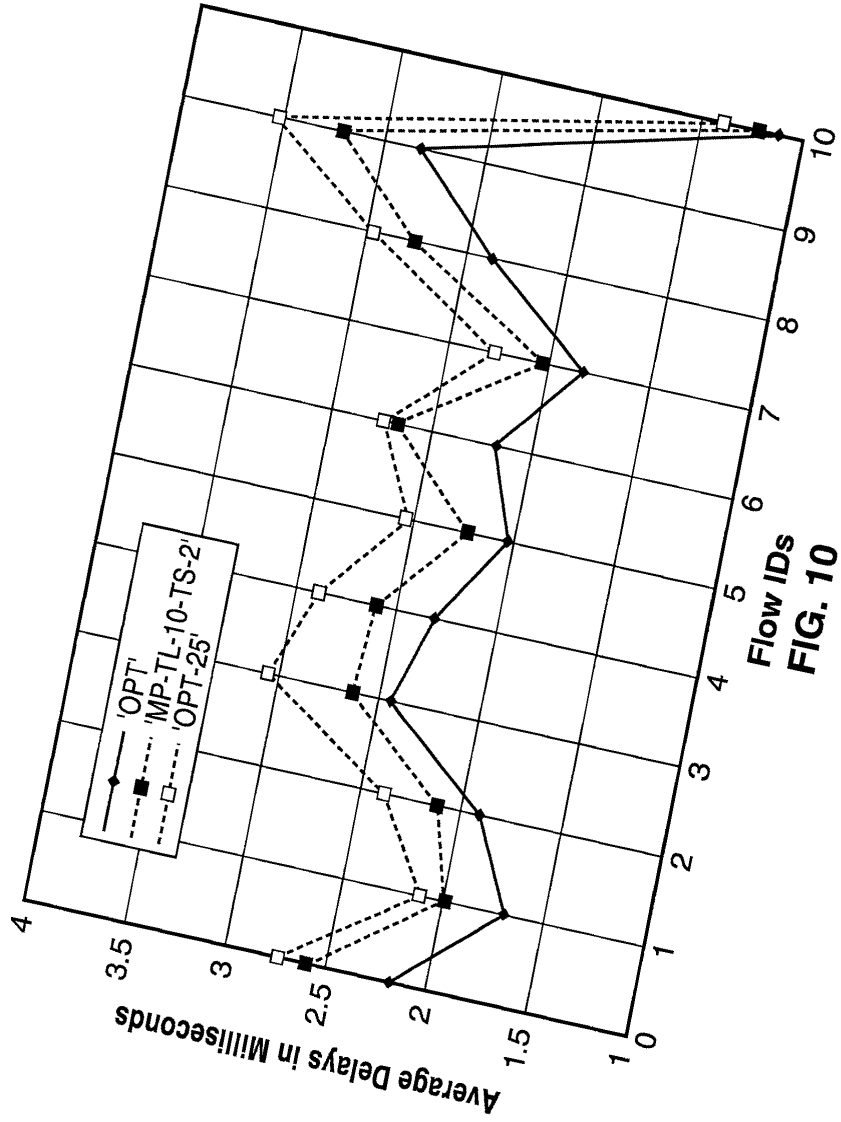


FIG. 10

11/19

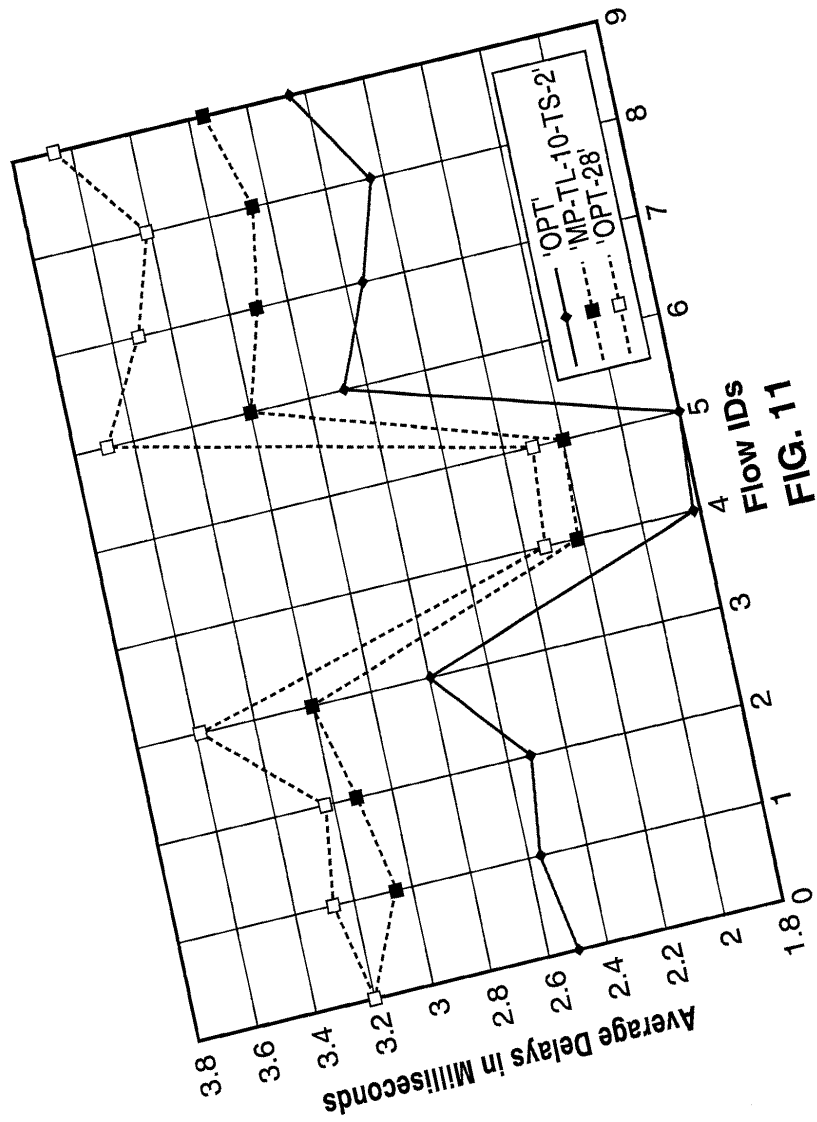
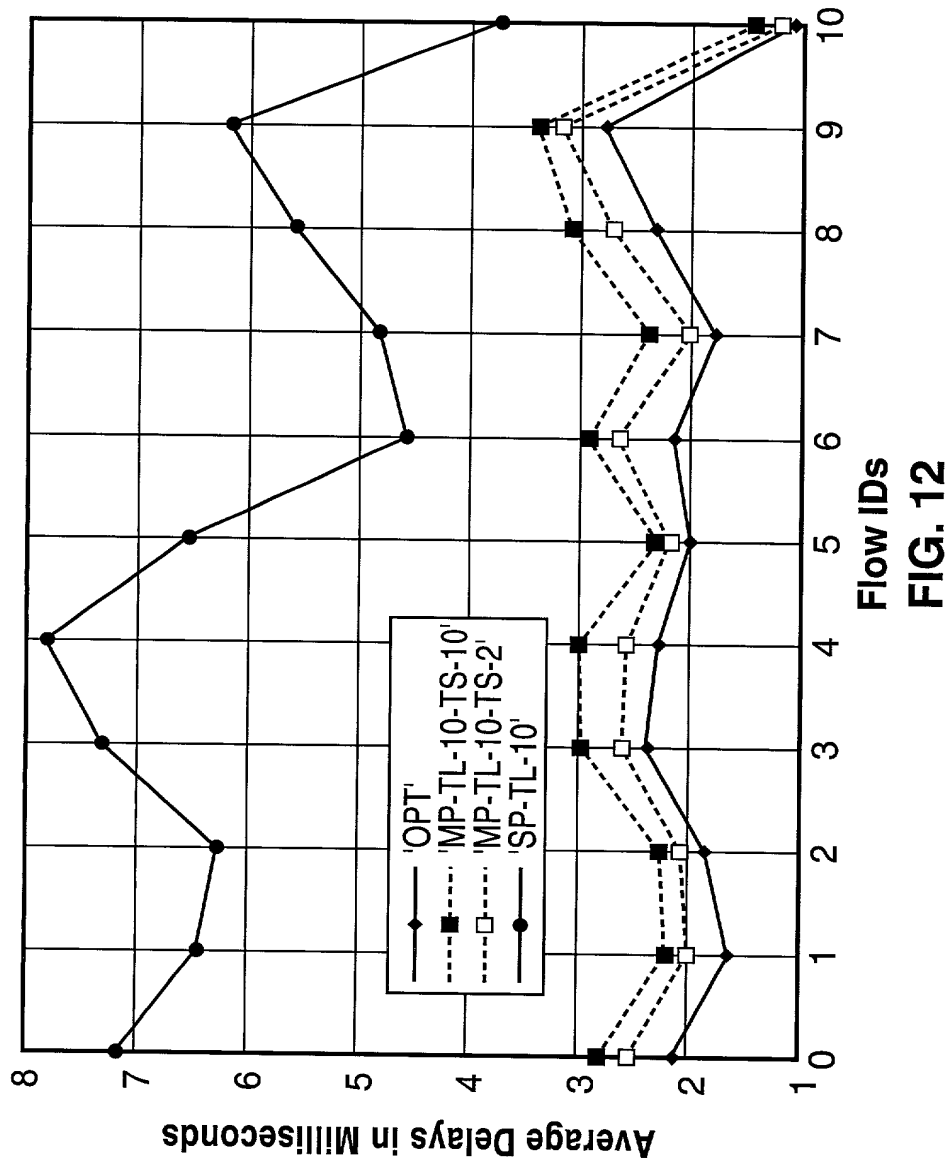


FIG. 11



Flow IDs
FIG. 12

1000 900 800 700 600 500 400 300 200 100 0

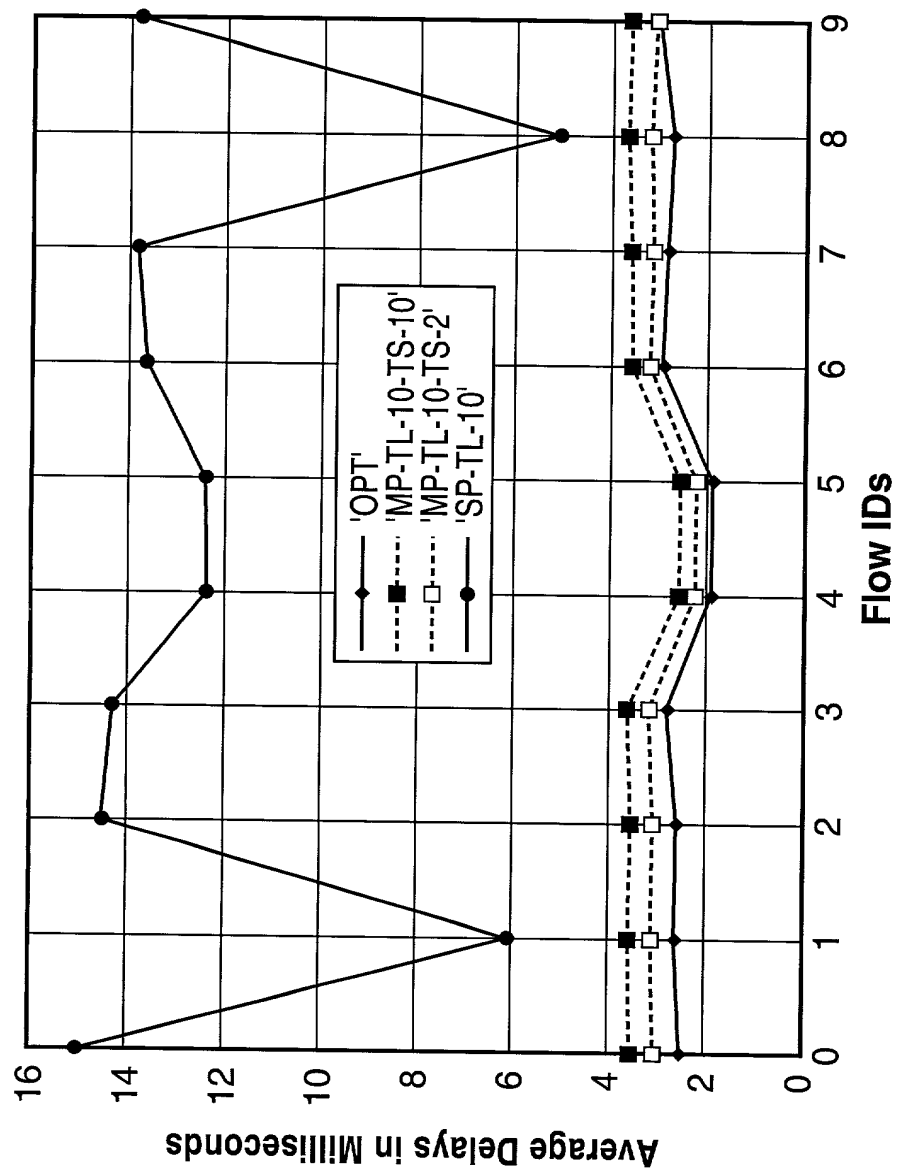
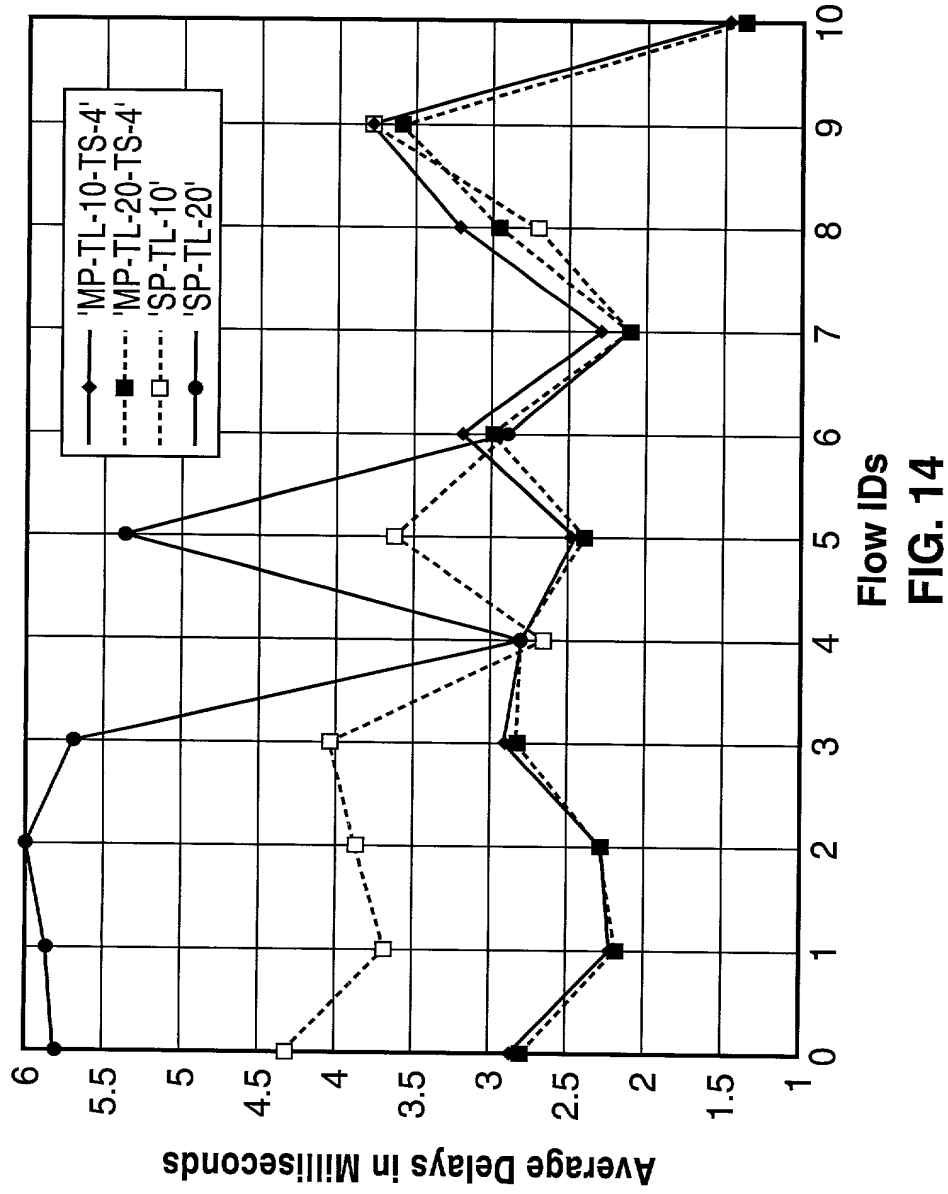
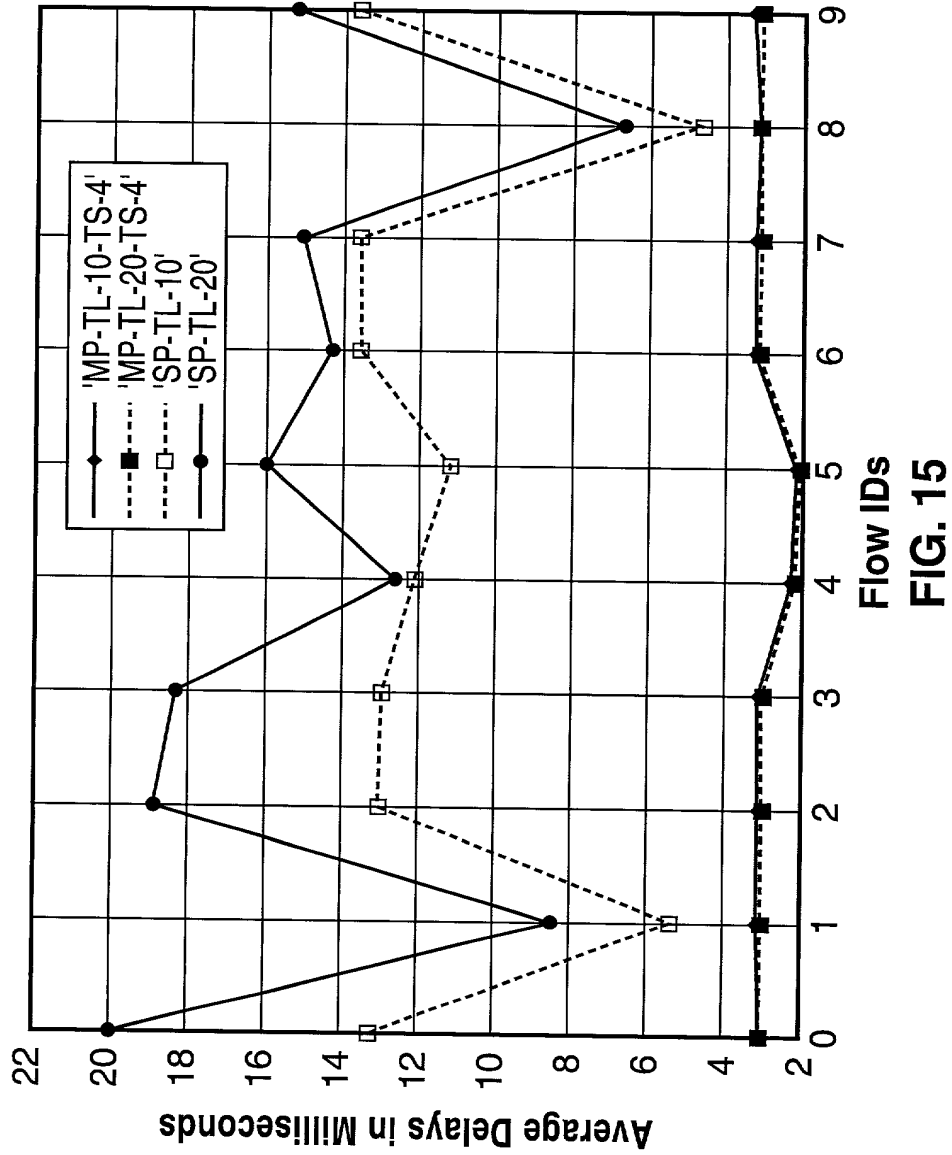


FIG. 13



Flow IDs
FIG. 14



Flow IDs
FIG. 15

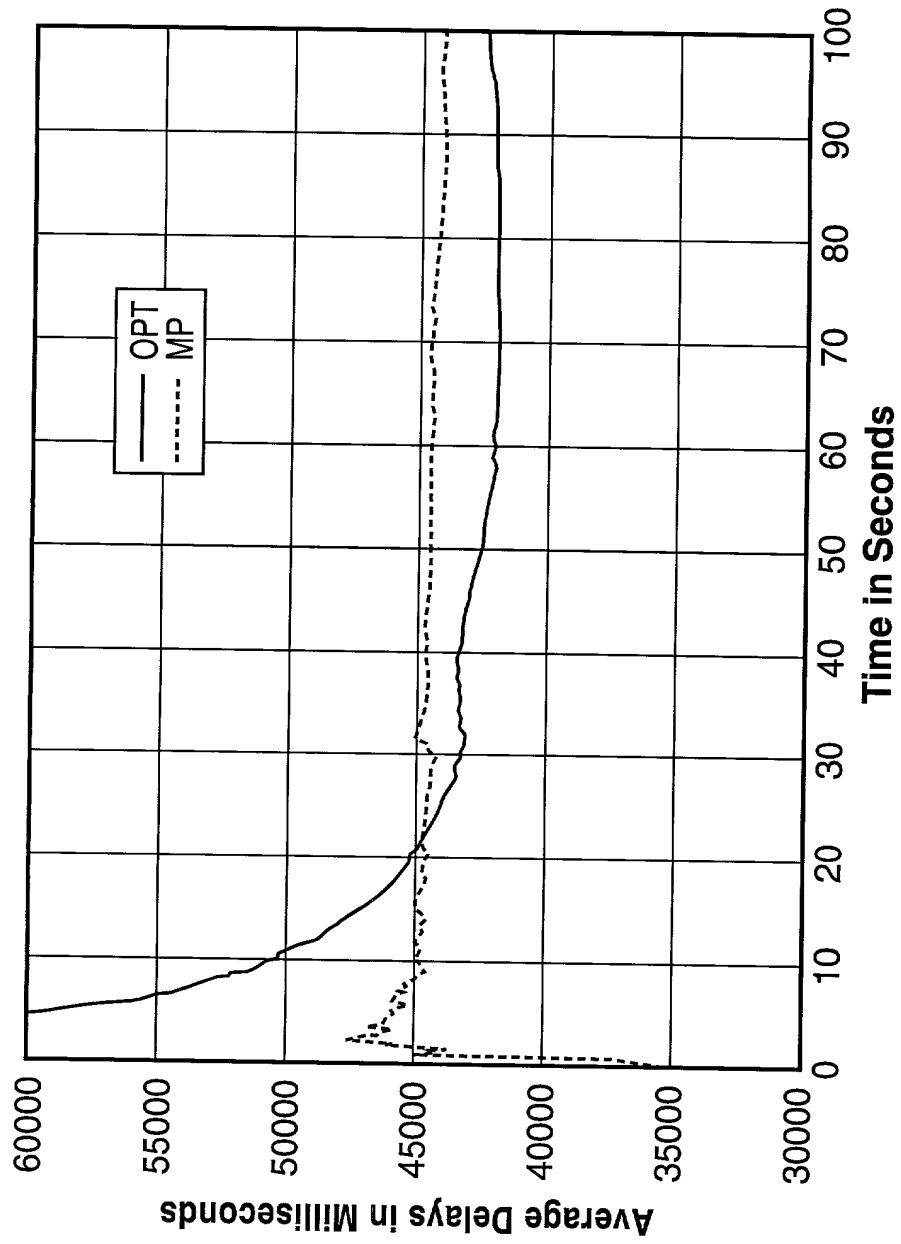


FIG. 16

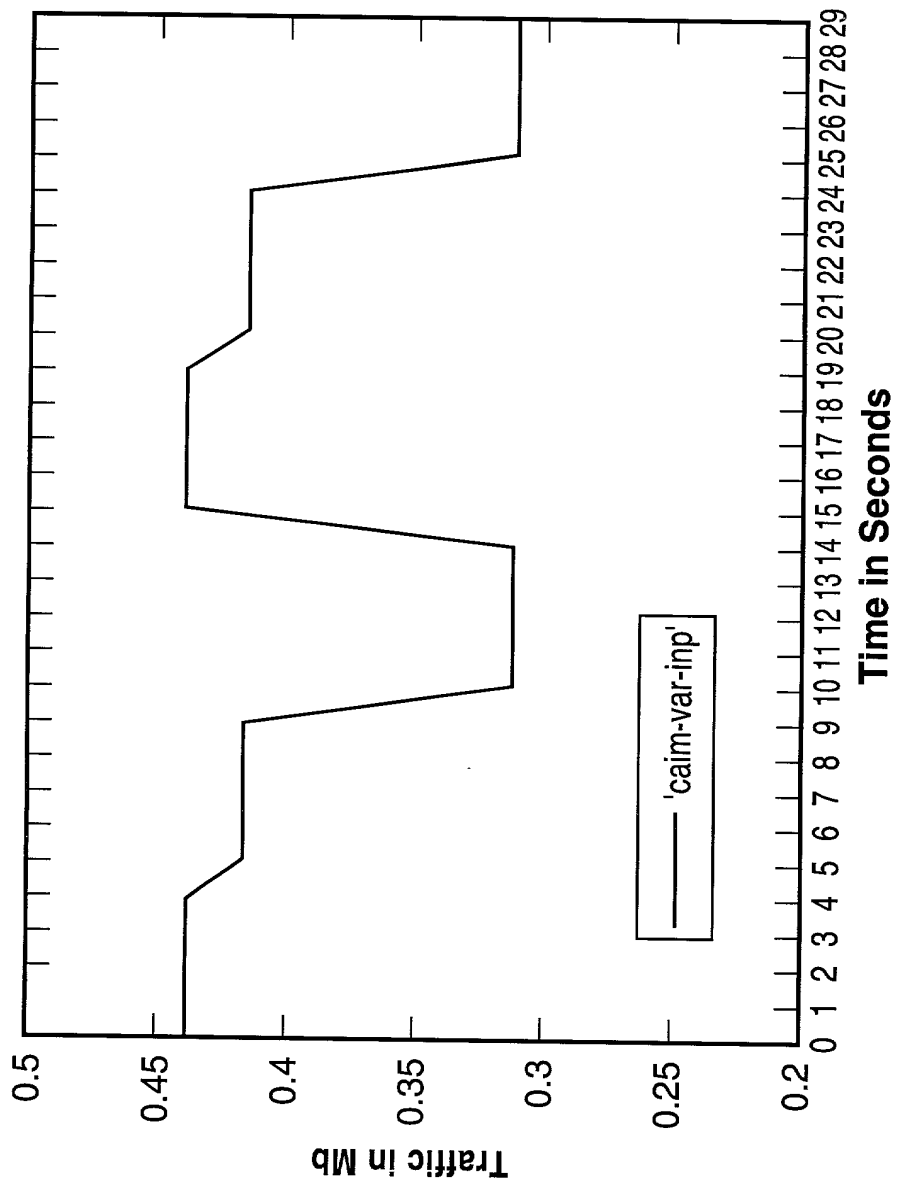


FIG. 17

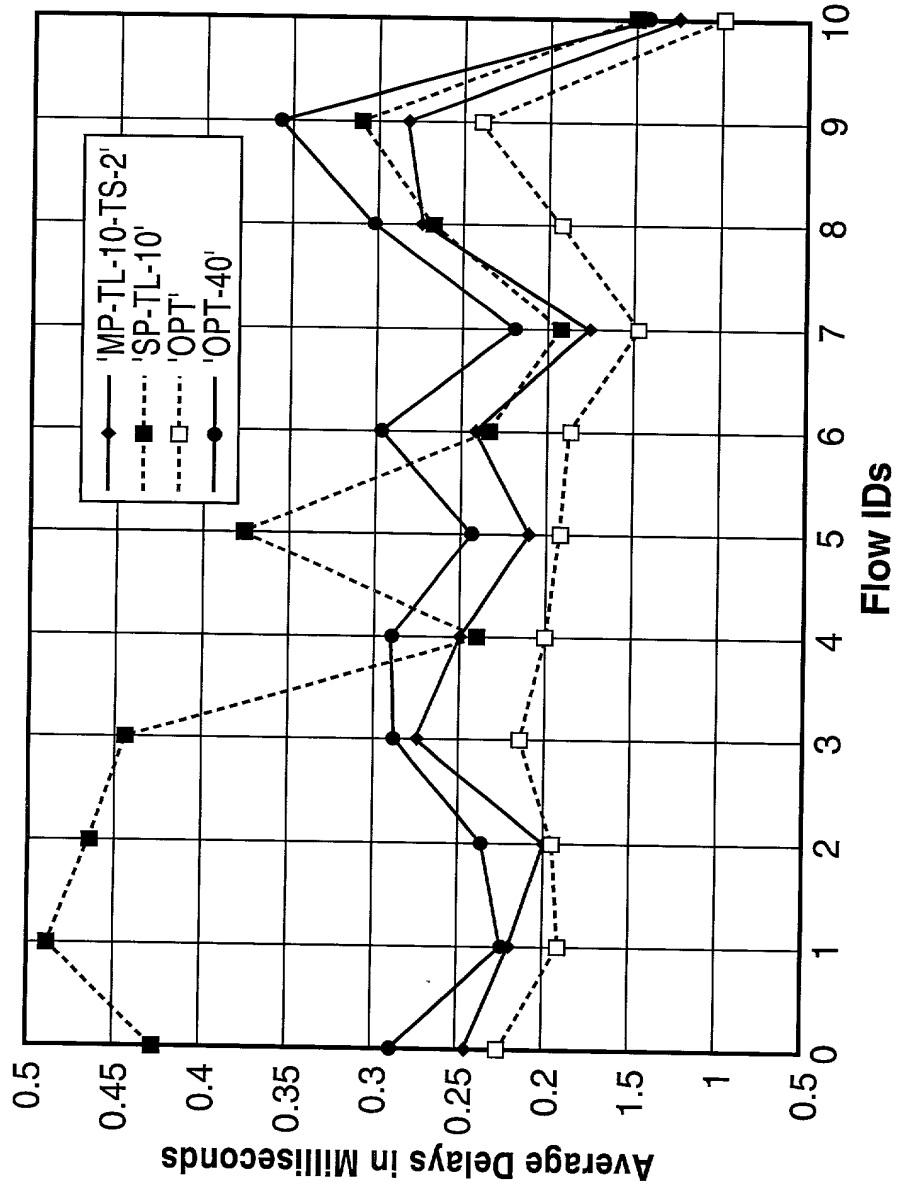


FIG. 18

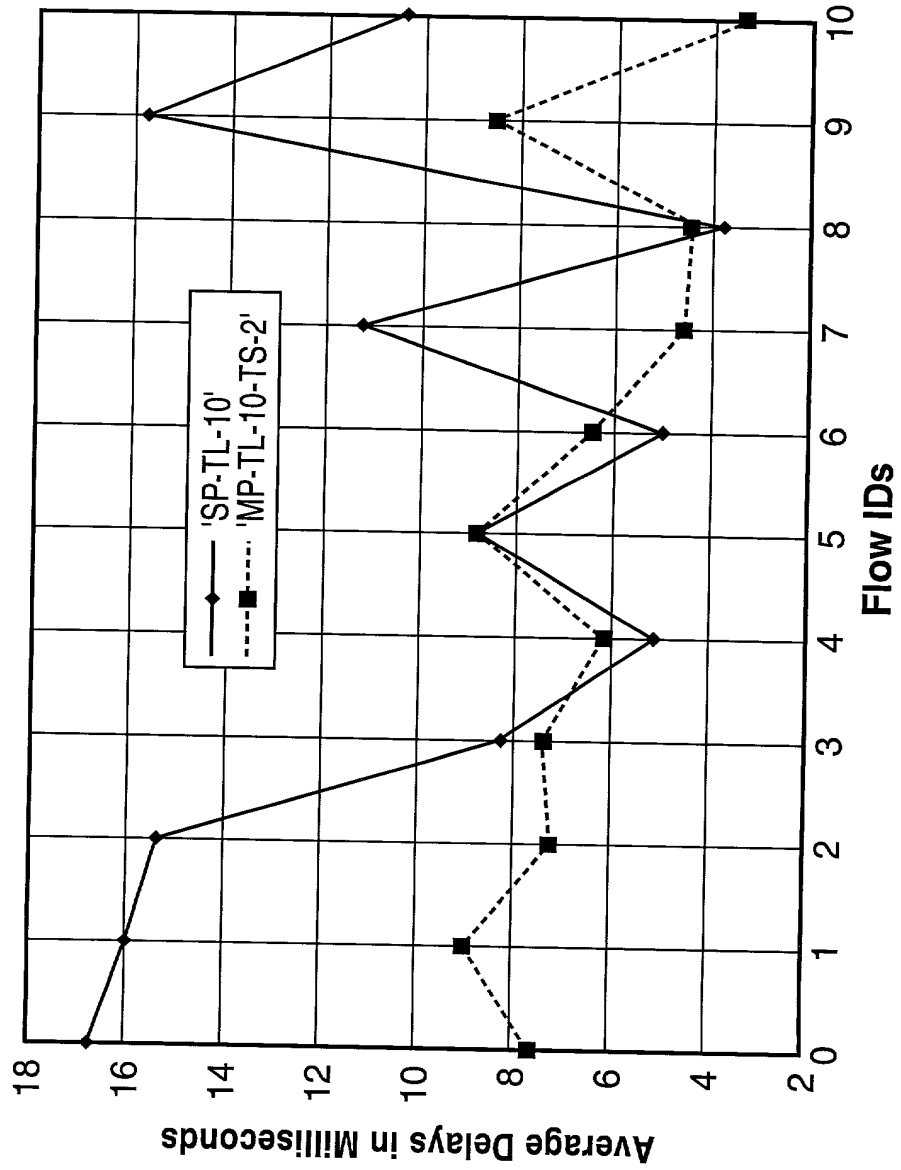


FIG. 19